Data Biased Robust Counter Strategies

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Introduction







• Computer Poker Research Group

- Created Polaris the world's strongest program for playing Heads-Up Limit Texas Hold'em Poker
- July 2008: Went to Las Vegas, played against six poker pros, won the 2nd Man-Machine Poker Championship
- Won several events in the 2008 AAAI Computer Poker Competition
- Research goals:
 - Solve very large extensive form games
 - Learn to model and exploit opponent's strategy

In this talk, we present a technique for dealing with three types of model uncertainty:

- The opponent / environment changes after we model it
- The model is more accurate in some areas than others
- The model's prior beliefs are very inaccurate

• Our domain: 2-player Limit Texas Hold'em Poker

- Zero-Sum Extensive form game
- Repeated game (Hundreds or thousands of short games)
- Hidden information (Can't see opponent's cards)
- Stochastic elements (Cards are dealt randomly)
- Goal: Win as much money as possible
- RL interpretation:
 - POMDP (when opponent's strategy is static)
 - Some properties of world are known
 - Probability distribution at chance nodes
 - Don't know exactly what state you are in (because of opponent's cards)
 - Transition probabilities at opponent choice nodes are unknown
 - Payoffs at terminal nodes are unknown

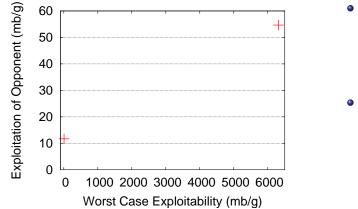
Types of strategies

- There are lots of ways to play games like poker. Two are well known:
 - Nash Equilibrium
 - Minimizes worst-case performance
 - Doesn't try to exploit opponent's mistakes
 - Best Response
 - Maximizes performance against a specific static opponent
 - Doesn't try to minimize worst-case performance
 - Problem: requires the opponent's strategy
- Goals:
 - Observe the opponent, build a model, and use it instead of the opponent's strategy

- Bound worst-case performance
 - Model could be inaccurate
 - Opponent could change

Types of Strategies

Performance against a static opponent, in millibets per game



 Game Theory: Nash equilibrium. Low exploitiveness, low exploitability
Decision Theory:

Best response.

exploitiveness,

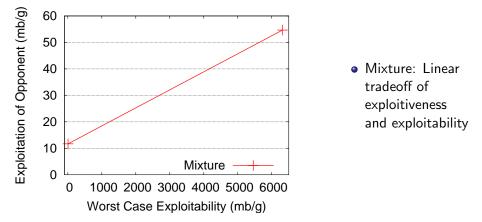
high exploitability

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Types of Strategies

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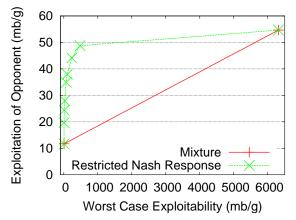


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Types of Strategies

Performance against a static opponent, in millibets per game



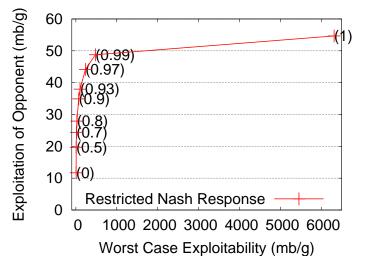
 Restricted Nash Response: Much better than linear tradeoff

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- Restricted Nash Response
 - Proposed by Johanson, Zinkevich and Bowling (Computing robust counter-strategies, NIPS 2007)
- Choose a value *p* and play an unusual game:
 - With probability *p*, opponent is forced to play according to a static strategy
 - With probability 1 p, opponent is free to play as they like
- p = 1: Best response
- p = 0: Nash equilibrium
- 0 robust to any opponent!
- This provably generates the best possible counter-strategies to the opponent

Restricted Nash Response

Performance against model of Orange



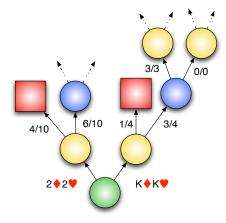
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• Observe the opponent, build a model, and use it instead of the opponent's strategy

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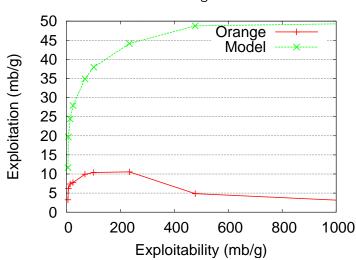
- Bound worst-case performance
 - Model could be inaccurate
 - Opponent could change

Frequentist Opponent Models



- Observe 100,000 to 1 million games played by the opponent
- Do frequency counts on actions taken at information sets
- Model assumes opponent takes actions with observed frequencies
- Need a default policy when there are no observations
 - Poker: Always-Call

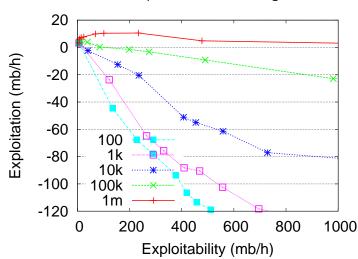
Problems with Restricted Nash Response



Problem 1: Overfitting to the model

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Problems with Restricted Nash Response



Problem 2: Requires a lot of training data

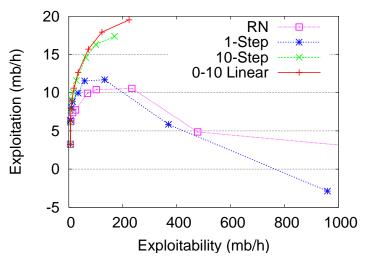
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- Restricted Nash Response had two problems:
 - Model wasn't accurate in states we never observed
 - Model was more accurate in some states than in others
- We need a new approach. We'd like to only use the model wherever we have reason to trust it
- New approach: use model's accuracy as part of the restricted game

- Lets set up another restricted game. Instead of one p value for the whole tree, we'll set one p value for each choice node, p(i)
- More observations \rightarrow more confidence in the model \rightarrow higher p(i)
- Set a maximum p(i) value, P_{\max} , that we vary to produce a range of strategies

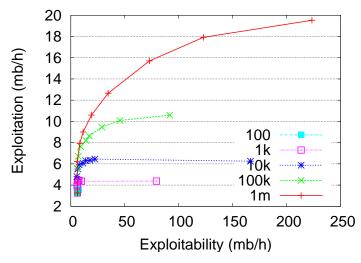
- Three examples:
 - 1-Step: p(i) = 0 if 0 observations, $p(i) = P_{\max}$ otherwise
 - 10-Step: p(i) = 0 if less than 10 observations, $p(i) = P_{\max}$ otherwise
 - 0-10 Linear: p(i) = 0 if 0 observations, $p(i) = P_{\max}$ if 10 or more, and p(i) grows linearly in between

 By setting p(i) = 0 in unobserved states, our prior is that the opponent will play as strongly as possible RNR and several DBR curves:



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- Data Biased Response technique:
 - Generate a range of strategies, trading off exploitation and worst-case performance

- Take advantage of observed information
- Avoid overfitting to parts of the model we suspect are inaccurate

- Extend to single-player domains
 - Can overfitting be reduced by assuming a slightly adversarial environment in unobserved / underobserved areas?

• More rigorous method for setting *p* from the observations